

# A Penalty-Interior-Point Algorithm for Nonlinear Optimization

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# Outline

Introduction

Algorithmic Framework

Parameter Updates

Numerical Experiments

Summary and Future Work

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## Introduction

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## Large-scale optimization

Consider the optimization problem:

$$(OP) := \begin{array}{l} \min_x f(x) \\ \text{s.t. } c(x) \leq 0 \end{array}$$

For large-scale instances:

- ▶ Linear or quadratic optimization subproblems are expensive. (Linear systems OK.)
- ▶ The constraints may be difficult to satisfy.
- ▶ The constraints may be (locally) infeasible; i.e., the algorithm should solve:

$$(FP) := \min_x v(x) := \sum_{i \in \mathcal{I}} \max\{c^i(x), 0\}$$

## Penalty methods

Unconstrained techniques can be used if we solve:

$$\min_x \rho f(x) + v(x)$$

Similarly, we can solve a regularized form of (OP):

$$\begin{aligned} \text{(PP)} := \quad & \min_{x,s} \rho f(x) + \sum_{i \in \mathcal{I}} s^i \\ & \text{s.t. } c(x) - s \leq 0, \quad s \geq 0 \end{aligned}$$

- ▶ Unconstrained techniques may fail or be slow if  $f$  is unbounded below; performance depends greatly on the form of  $v$ .
- ▶ Solving (PP) commonly requires the solution of linear or quadratic subproblems.
- ▶ Either way, **updating the penalty parameter is a challenge.**

## Interior-point methods

Large-scale problems are often solved efficiently through interior-point subproblems:

$$\text{(IP)} := \begin{array}{l} \min_{x,r} f(x) - \mu \sum_{i \in \mathcal{I}} \ln r^i \\ \text{s.t. } c(x) + r = 0 \quad (\text{with } r > 0) \end{array}$$

- ▶ Lacks constraint regularization as in a penalty method.
- ▶ Similar to before, **updating the interior-point parameter is a challenge.**

## Penalty-interior-point methods(?)

Can penalty and interior-point ideas be combined to create a **practical** algorithm?

- ▶ Regularization through penalties is an attractive feature.
- ▶ Search direction computations via linear system solves is nice for large problems.

However, there are significant challenges:

- ▶ Penalty methods want the algorithm to be free to violate constraints.
- ▶ Interior-point methods want the algorithm to remain feasible.
- ▶ **Juggling “conflicting” parameters is a major challenge.**

## Literature

Previous work with similar motivations:

- ▶ Jittorntrum and Osborne (1980)
- ▶ Polyak (1982, 1992, 2008)
- ▶ Breitfeld and Shanno (1994, 1996)
- ▶ Goldfarb, Polyak, Scheinberg, and Yuzefovich (1999)
- ▶ **Gould, Orban, and Toint (2003)**
- ▶ Chen and Goldfarb (2006, 2006)
- ▶ Benson, Sen, and Shanno (2008)

However, in my view, none of these papers focus enough on parameter updates.



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## Penalty-interior-point subproblem

Recall:

$$\begin{aligned}
 \text{(PP)} := & \boxed{\begin{array}{l} \min_{x,s} \rho f(x) + \sum_{i \in \mathcal{I}} s^i \\ \text{s.t. } c(x) - s \leq 0, s \geq 0 \end{array}} & \text{(IP)} := & \boxed{\begin{array}{l} \min_{x,r} f(x) - \mu \sum_{i \in \mathcal{I}} \ln r^i \\ \text{s.t. } c(x) + r = 0 \text{ (with } r > 0) \end{array}}
 \end{aligned}$$

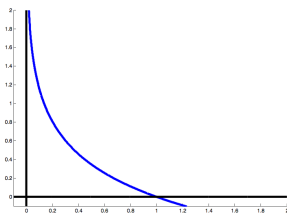
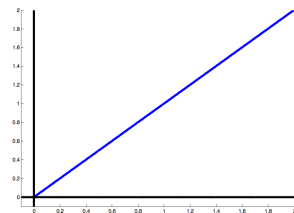
Applying an interior-point reformulation to (PP), we can obtain:

$$\text{(PIP)} := \boxed{\begin{array}{l} \min_{x,r,s} \rho f(x) - \mu \sum_{i \in \mathcal{I}} (\ln r^i + \ln s^i) + \sum_{i \in \mathcal{I}} s^i \\ \text{s.t. } c(x) + r - s = 0 \text{ (with } r, s > 0) \end{array}}$$

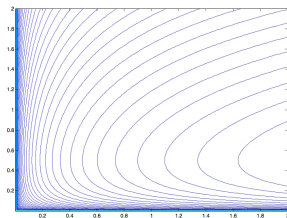
- ▶ (PIP) satisfies MFCQ (it is a reformulation of (PP), which also satisfies it).
- ▶  $\mu \rightarrow 0$  and  $\rho \rightarrow \bar{\rho} > 0$  to obtain a solution to (OP).
- ▶  $\mu \rightarrow 0$  and  $\rho \rightarrow 0$  to obtain a solution to (FP).

## Visualizing the penalty-interior-point objective

- ▶ Objective function terms for  $s^i$  in (PP) and  $r^i$  in (IP):



- ▶ Objective function term for  $(r^i, s^i)$  in (PIP):



## Algorithm outline

for  $k = 0, 1, 2, \dots$

- ▶ Reset the slack variables.
- ▶ **Update the parameters.**
- ▶ Compute a search direction.
- ▶ Perform a line search.

## Slack reset

Through the slack variables, we have added many degrees of freedom to the problem!

- ▶ However, for a fixed  $x_k$ , (PIP) reduces to

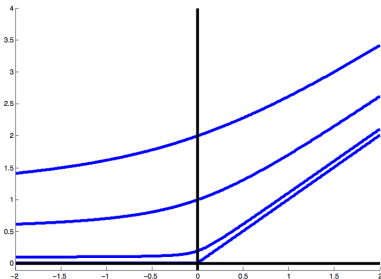
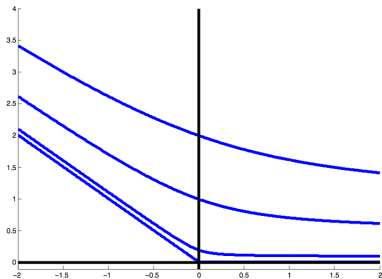
$$\begin{aligned} \min_{r,s} \quad & -\mu \sum_{i \in \mathcal{I}} (\ln r^i + \ln s^i) + \sum_{i \in \mathcal{I}} s^i \\ \text{s.t.} \quad & c(x_k) + r - s = 0 \quad (\text{with } r, s > 0) \end{aligned}$$

- ▶ This problem is convex and separable, and has the unique solution:

$$\begin{aligned} r_k^i &= r^i(x_k; \mu) := \mu - \frac{1}{2} c^i(x_k) + \frac{1}{2} \sqrt{c^i(x_k)^2 + 4\mu^2} \\ \text{and } s_k^i &= s^i(x_k; \mu) := \mu + \frac{1}{2} c^i(x_k) + \frac{1}{2} \sqrt{c^i(x_k)^2 + 4\mu^2}. \end{aligned}$$

## Visualizing the slack reset

Slack variables  $r$  and  $s$ , respectively, as functions of  $\mu$  and  $c(x_k)$ :



## Search direction calculation

A Newton iteration for the optimality conditions of (PIP) involves:

$$\begin{bmatrix} H_k & 0 & 0 & \nabla c(x_k) \\ 0 & \Omega_k & 0 & I \\ 0 & 0 & \Gamma_k & -I \\ \nabla c(x_k)^T & I & -I & 0 \end{bmatrix} \begin{bmatrix} \Delta x_k \\ \Delta r_k \\ \Delta s_k \\ \Delta \lambda_k \end{bmatrix} = - \begin{bmatrix} \rho \nabla f(x_k) + \nabla c(x_k) \lambda_k \\ \lambda_k - \mu R_k^{-1} e \\ e - \lambda_k - \mu S_k^{-1} e \\ c(x_k) + r_k - s_k \end{bmatrix}$$

(Why do we still have  $(\Delta r_k, \Delta s_k)$  if we eliminated  $(r_k, s_k)$ ?)

## Merit function

- ▶ Recall that the objective of (PIP) is given by

$$\phi(x, r, s; \rho, \mu) := \rho f(x) - \mu \sum_{i \in \mathcal{I}} (\ln r^i + \ln s^i) + \sum_{i \in \mathcal{I}} s^i.$$

- ▶ A standard type of merit function or filter for (PIP) would involve  $\phi$  and a measure of violation of the constraints  $c(x) + r - s = 0$ .
- ▶ However, the slack reset allows us to use the merit function

$$\tilde{\phi}(x; \rho, \mu) := \rho f(x) - \mu \sum_{i \in \mathcal{I}} (\ln r^i(x; \mu) + \ln s^i(x; \mu)) + \sum_{i \in \mathcal{I}} s^i(x; \mu).$$

### Lemma

Let  $r_k = r(x_k; \mu)$  and  $s_k = s(x_k; \mu)$ . Then, the computed search direction  $\Delta x_k$  yielded by the Newton system is a descent direction for  $\tilde{\phi}(x; \rho, \mu)$  at  $x = x_k$ .



## Line search

For a given search direction  $(\Delta x_k, \Delta \lambda_k)$ , we:

- ▶ backtrack to find  $\alpha_k \in (0, 1]$  satisfying the fraction-to-the-boundary rules

$$r(x_k + \alpha \Delta x_k; \mu) \geq \tau r_k$$

and

$$s(x_k + \alpha \Delta x_k; \mu) \geq \tau s_k$$

and the sufficient decrease condition

$$\tilde{\phi}(x_k + \alpha_k \Delta x_k; \rho, \mu) \leq \tilde{\phi}(x_k; \rho, \mu) + \eta \alpha_k \nabla \tilde{\phi}(x_k; \rho, \mu)^T \Delta x_k.$$

- ▶ compute the largest  $\beta_k \in (0, 1]$  satisfying the fraction-to-the-boundary rule

$$\lambda_k + \beta \Delta \lambda_k \in [\tau \lambda_k, e - \tau(e - \lambda_k)].$$

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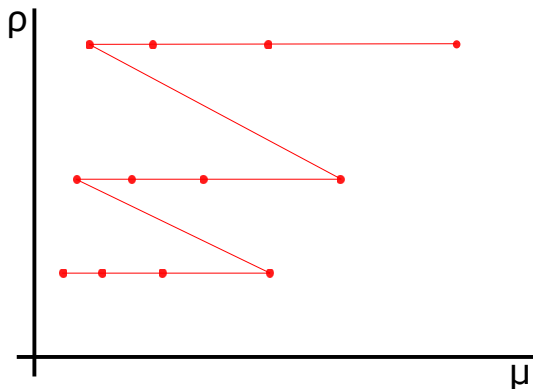
## Conservative strategies

A simple, conservative strategy may be the following:

- ▶ Step 1: Fix  $\rho$  and solve (PIP) for  $\mu \rightarrow 0$ .
- ▶ Step 2: If we are infeasible, decrease  $\rho$  and go to step 1.

This strategy, or ones that are equally as conservative, are the type that have been implemented in many other penalty-interior-point algorithms.

## Visualizing a conservative strategy



- ▶ Each “dot” may require at least a few iterations.
- ▶ Each “row” may require the computational effort of an entire interior-point run!
- ▶ For an infeasible problem, we need both  $\rho$  and  $\mu$  to reduce to (near) zero.

## Search direction calculation

Recall the Newton system:

$$\begin{bmatrix} H_k & 0 & 0 & \nabla c(x_k) \\ 0 & \Omega_k & 0 & I \\ 0 & 0 & \Gamma_k & -I \\ \nabla c(x_k)^T & I & -I & 0 \end{bmatrix} \begin{bmatrix} \Delta x_k \\ \Delta r_k \\ \Delta s_k \\ \Delta \lambda_k \end{bmatrix} = - \begin{bmatrix} \rho \nabla f(x_k) + \nabla c(x_k) \lambda_k \\ \lambda_k - \mu R_k^{-1} e \\ e - \lambda_k - \mu S_k^{-1} e \\ 0 \end{bmatrix}$$

- ▶  $\rho$  and  $\mu$  may be embedded in  $H_k$ ,  $\Omega_k$ , and  $\Gamma_k$ , but...
- ▶ Holding these matrices fixed for iteration  $k$  means we have a system of the form:

$$M \Delta z_k^{\rho, \mu} = \rho \begin{bmatrix} -\nabla f(x_k) \\ 0 \\ 0 \\ 0 \end{bmatrix} + \mu \begin{bmatrix} 0 \\ R_k^{-1} e \\ S_k^{-1} e \\ 0 \end{bmatrix} + \begin{bmatrix} -\nabla c(x_k)^T \lambda_k \\ -\lambda_k \\ -e + \lambda_k \\ 0 \end{bmatrix}.$$

- ▶ Thus, the solution for **all** pairs  $(\rho, \mu)$  can be obtained with only one factorization.

## Update criteria

- ▶ It is computationally practical to vary  $\rho$  and  $\mu$  on the right-hand side.
- ▶ Ok, but what criteria should we use for choosing these values?
- ▶ In a **penalty method**, decreasing  $\rho$  places more emphasis on the constraints.
- ▶ In an **interior-point method**, decreasing  $\mu$  places less emphasis on centrality.
- ▶ However, in a **penalty-interior-point method**, everything gets jumbled!

In short, we update:

- ▶  $\rho$  to ensure some level of progress toward solving the **primal feasibility** problem;
- ▶  $\mu$  to attempt to satisfy **dual feasibility** and **complementarity**.

## A basis for comparison

Let  $z := (x, r, s)$ .

- ▶ We have two views of the penalty-interior-point objective:

$$\phi(z; \rho, \mu) = \rho f(x) - \mu \sum_{i \in \mathcal{I}} (\ln r^i + \ln s^i) + \sum_{i \in \mathcal{I}} s^i;$$

$$\tilde{\phi}(x; \rho, \mu) = \rho f(x) - \mu \sum_{i \in \mathcal{I}} (\ln r^i(x; \mu) + \ln s^i(x; \mu)) + \sum_{i \in \mathcal{I}} s^i(x; \mu).$$

- ▶ Thus, we have two corresponding linear models:

$$l(\Delta z; \rho, \mu, z) := \phi(z; \rho, \mu) + \nabla \phi(z; \rho, \mu)^T \Delta z;$$

$$\tilde{l}(\Delta x; \rho, \mu, x) := \tilde{\phi}(x; \rho, \mu) + \nabla \tilde{\phi}(x; \rho, \mu)^T \Delta x.$$

- ▶ For the  $\mu$  used in the slack reset, the models coincide, but not otherwise.
- ▶ It is easily seen in the direction computation that

$$\Delta l(\Delta z; \rho, \mu, z) := l(0; \rho, \mu, z) - l(\Delta z; \rho, \mu, z) > 0,$$

but it is a reduction in  $\tilde{l}(\cdot; \rho, \mu, x)$  that we want to guarantee!

## Updating $\rho$ : ensuring progress in feasibility

Let  $\Delta z_k^{\rho, \mu}$  be the direction computed for given  $(\rho, \mu)$ .

- ▶ If  $x_k$  is feasible, then choose largest  $\rho$  such that for some  $\mu$ :

$$\Delta \tilde{q}(\Delta x_k^{\rho, \mu}; \rho, \mu, x_k) > 0.$$

(Here,  $\tilde{q}(\Delta x; \rho, \mu, x)$  is a quadratic model of  $\tilde{\phi}(x; \rho, \mu)$ .)

- ▶ If  $x_k$  is infeasible, then choose largest  $\rho$  such that for some  $\mu$ :

$$\begin{aligned} \Delta \tilde{l}(\Delta x_k^{\rho, \mu}; \rho, \mu, x_k) &\geq \epsilon_1 \Delta l(\Delta z_k^{0, \mu}; 0, \mu, z_k), \quad \epsilon_1 \in (0, 1); \\ \Delta \tilde{q}(\Delta x_k^{\rho, \mu}; \rho, \mu, x_k) &\geq \epsilon_2 \Delta l(\Delta z_k^{0, \mu}; 0, \mu, z_k), \quad \epsilon_2 \in (0, 1). \end{aligned}$$



## Updating $\rho$ : promoting fast infeasibility detection

If  $x_k$  is infeasible, then  $\rho$  must satisfy

$$\rho \leq \left\| \begin{bmatrix} \nabla c(x_k) \lambda_k \\ R_k \lambda_k \\ S_k (e - \lambda_k) \end{bmatrix} \right\|^2$$

(Right-hand side only small in neighborhood of an infeasible stationary point.)

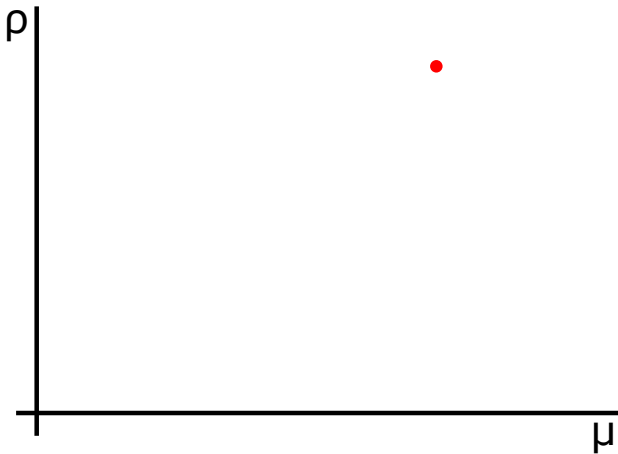
## Updating $\mu$ : minding dual feasibility and complementarity

Fixing  $\rho$ , now choose  $\mu$  so that the previous conditions still hold and

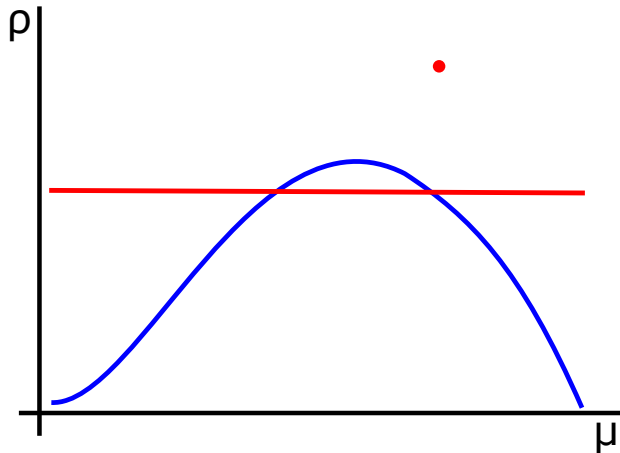
$$m(\Delta z, \Delta \lambda; \rho, \mu, z_k) := \left\| \begin{bmatrix} \rho \nabla f(x_k) + \nabla c(x_k)(\lambda_k + \Delta \lambda) \\ (R_k + \Delta R)(\lambda_k + \Delta \lambda) \\ (S_k + \Delta S)(e - \lambda_k - \Delta \lambda) \end{bmatrix} \right\|_{\infty}$$

is approximately minimized.

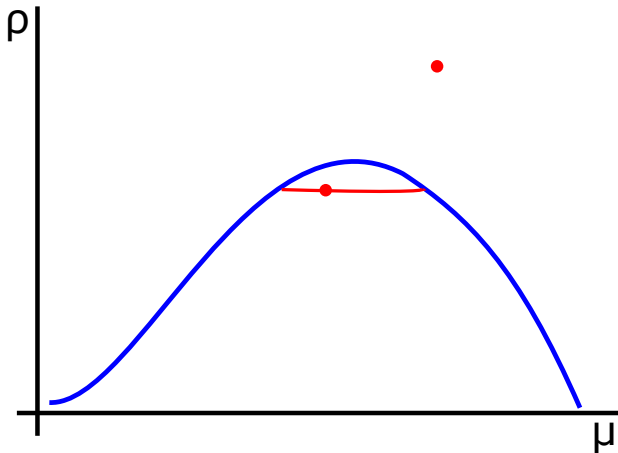
## Visualizing the aggressive strategy



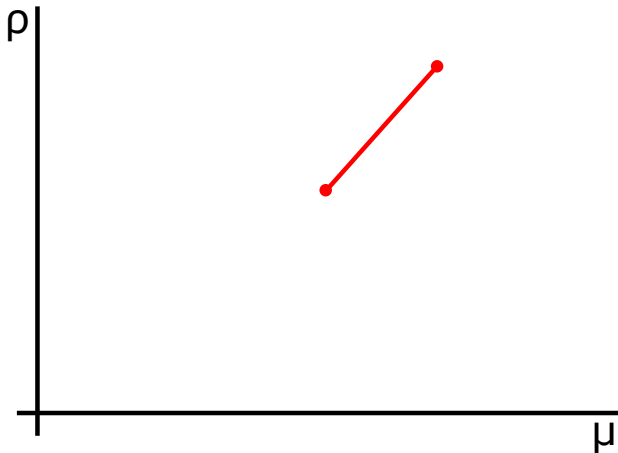
## Visualizing the aggressive strategy



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## Visualizing the aggressive strategy



All within one iteration!

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## Preliminary experimentation

- ▶ Penalty-Interior-Point ALgorithm (PIPAL)
- ▶ Compared iteration counts for PIPAL-c, PIPAL-a, and FMINCON.<sup>1</sup>
- ▶ CUTEr problems available in AMPL (385 in final set)
- ▶ Infeasible variants of the HS problems (93 in final set):

$$c'(x) = 0 \Rightarrow \{c'(x) = 0 \ \& \ c'(x) = 1\}$$

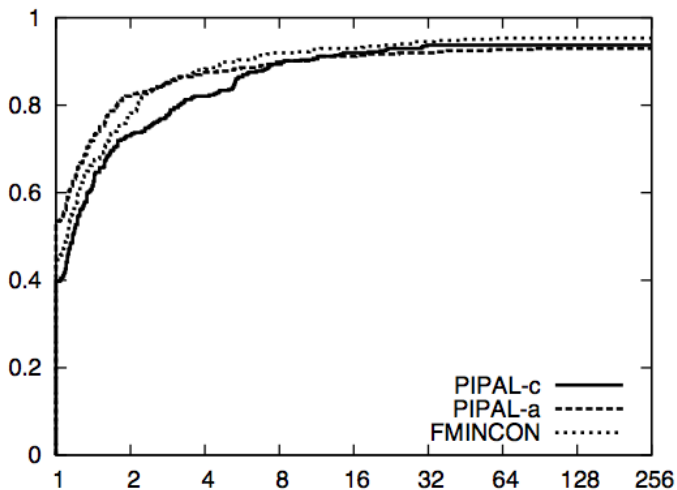
$$c'(x) \leq 0 \Rightarrow \{c'(x) \leq 0 \ \& \ c'(x) \geq 1\}$$

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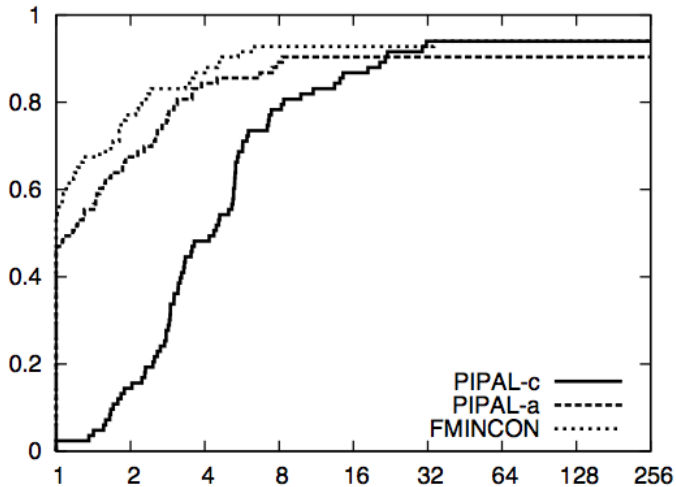
<sup>1</sup>See Waltz, Morales, Nocedal, and Orban (2006)



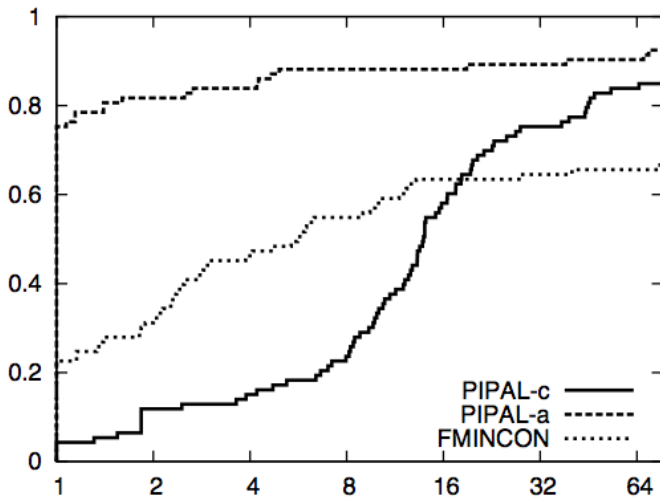
## Entire test set



## Problems requiring a penalty parameter decrease



## Infeasible problems



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## Summary

- ▶ Combined penalty and interior-point techniques into a single algorithm.
- ▶ Penalties effectively regularize the constraints.
- ▶ Interior-point strategies allowed directions to be computed via linear systems.
- ▶ Slack reset allowed us to reduce the degrees of freedom.
- ▶ **Proposed an aggressive updating scheme for the parameters.**
- ▶ Results are comparable to an interior-point method on most problems.
- ▶ Results are **much better** than an interior-point method **on infeasible problems.**

## Future work

- ▶ Advanced implementation is required — numerical issues are critical!
- ▶ Parameter update strategy for MPCCs — needs to be different?
- ▶ Specializations for MINLP — especially in terms of infeasibility detection?

## A new normal step?

- ▶ Our update for  $\rho$  involved comparing directions to that obtained with  $\rho = 0$ .
- ▶ That is, it is compared to an **interior-point step for the feasibility problem**:

$$\begin{aligned}
 \text{(FIP)} := & \min_{\Delta x, \Delta r, \Delta s} -\mu(R_k^{-1}\Delta r + S_k^{-1}\Delta s) + \sum_{i \in \mathcal{I}} \Delta s^i \\
 & \text{s.t. } (c(x_k) + \nabla c(x_k)^T \Delta x) + (r_k + \Delta r) - (s_k + \Delta s) = 0
 \end{aligned}$$

- ▶ In contrast, step decomposition methods attempt to solve:

$$\begin{aligned}
 \text{(NP)} := & \min_{\Delta x, \Delta s} \|c(x_k) + \nabla c(x_k)^T \Delta x - (s_k + S_k(S_k^{-1}\Delta s))\|^2 \\
 & \text{s.t. } \|\Delta x\|^2 + \|S_k^{-1}\Delta s\|^2 \leq \delta_k^2 \\
 & \quad S_k + \Delta s \geq \tau s_k
 \end{aligned}$$

- ▶ Does PIPAL offer an alternative to the standard normal subproblem?