

# Primal heuristics in MIPs

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# Outline of Talk

- Introduction: **MIP** and **Primal Heuristic**
- Description of two **primal heuristics**
- Implementation
- Proposed Work

## Primal Heuristics in MIP (1)

- A typical MIP solver uses **Branch** and **Bound** (and **Cuts**)
- A **Linear Program** is solved at each node
- There are 3 possibilities
  1. Solution of LP is feasible to the original problem
  2. Solution of the LP is non integral and hence infeasible for the original problem
  3. LP is infeasible
- If solution is non-integral, there are two possibilities
  1.  $z_{LP} \geq$  best solution  $\Rightarrow$  **Fathom**
  2.  $z_{LP} \leq$  best solution  $\Rightarrow$  **sub-branching**

Hence a good **incumbent** is important.

## Primal Heuristics in MIP (2)

A good incumbent may be achieved by:

- **Fixing** variables and diving
- Searching **around the current LP solution** for an integral solution
- If an IP solution is found, searching for a **better solution** in the neighbourhood.

None of the approaches guarantee a good solution.

None of them guarantee good speed

Good heuristics have been limited to specific problem structures

## Our Objective

To implement **generalized** primal heuristics for an MIP solver

## Proposed Heuristics

1. Local Branching (Fischetti and Lodi - 2002)
2. RINS (Danna, Rothberg and Pape - 2004)

## MIP solvers

1. MINTO
2. SYMPHONY

# LOCAL BRANCHING

- A form of **soft-fixing**: give solver some freedom to fix variables e.g.

$$\Delta(x, \bar{x}) = \sum_{j \in \bar{S}} (1 - x_j) + \sum_{j \in B \setminus \bar{S}} x_j \leq k$$

where,

$$\bar{S} = \{j \in B : \bar{x}_j = 1\}$$

- Branching decision

$$\Delta(x, \bar{x}^1) \leq k$$

$$\Delta(x, \bar{x}^1) \geq k + 1$$

- Iterative procedure
- Additionally time-bounds and diversification could be used.

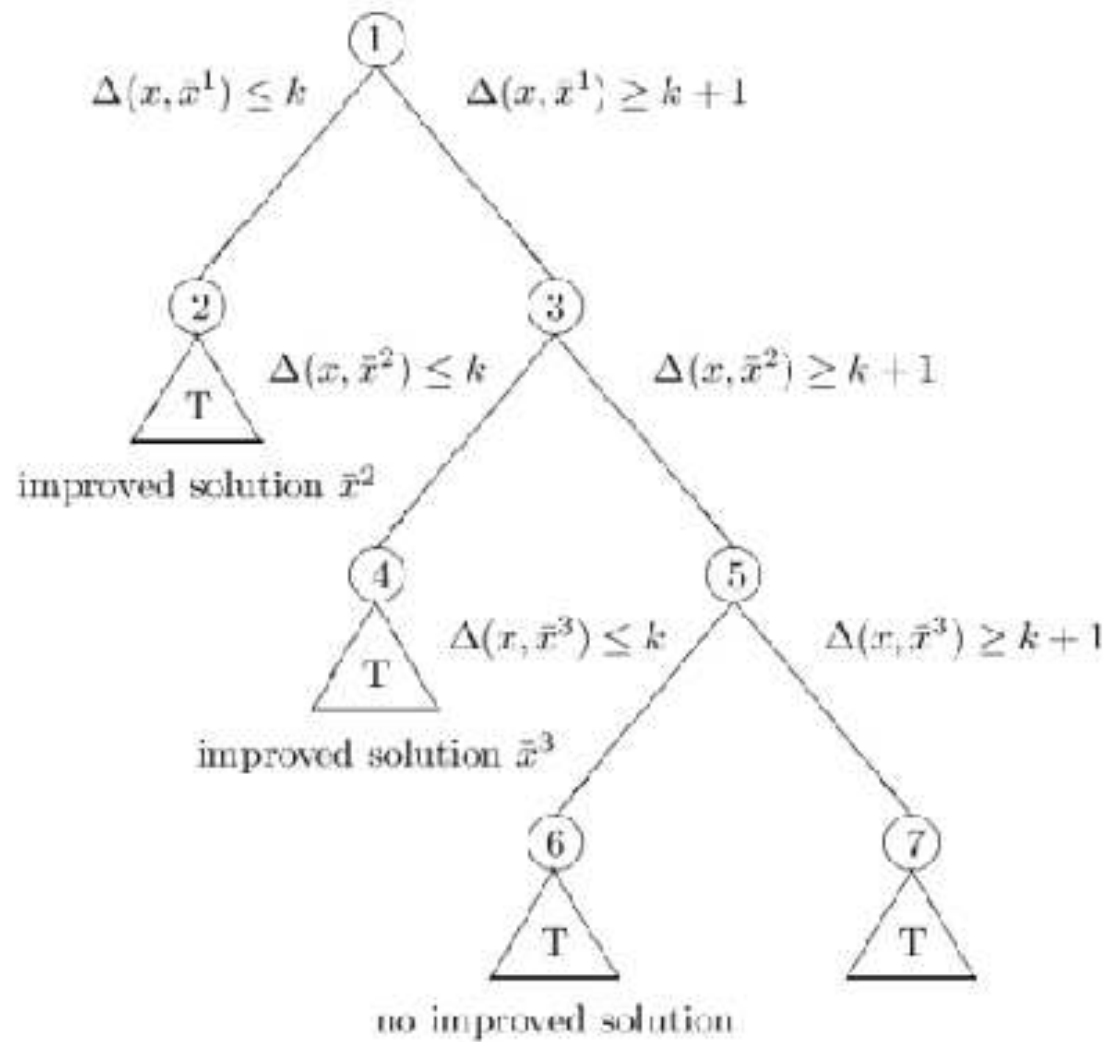


Figure 1: The basic local branching scheme.

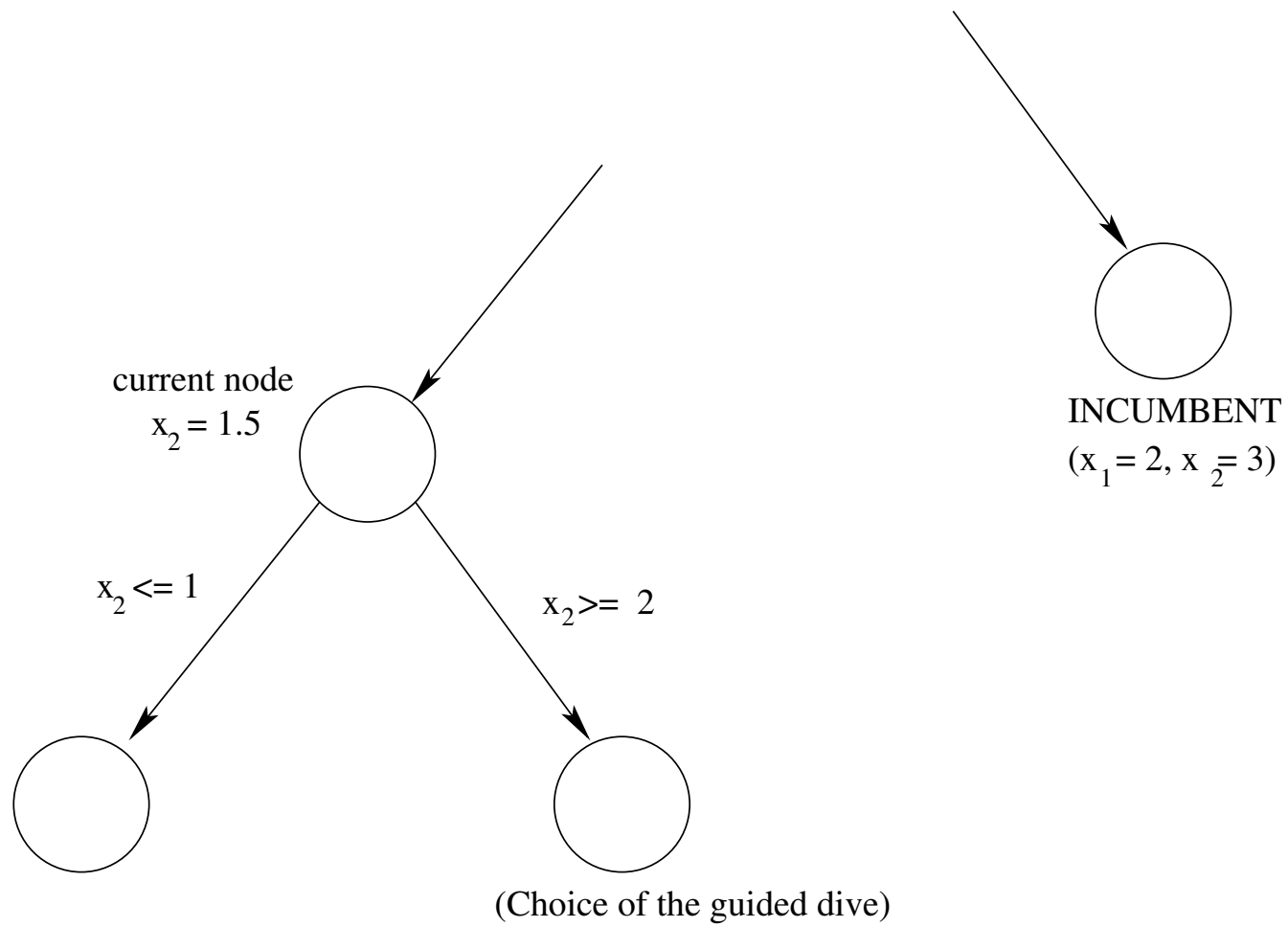
## Relaxation Induced Neighborhood Search (RINS)

- Assumes that an incumbent exists.
- Intuitively, a good solution should be around the **LP Optimal**. Also, it should have some closeness to the **incumbent**.
- Fix those variables which have the **same value** in the incumbent and the LP optimal.
- Form a new MIP after fixing these variables.
- Solve this simpler MIP.
- Optionally, put a limit on the number of nodes to be searched in the sub-MIP.



## Guided Dives

- Extends the idea of RINS while deciding **which branch** to choose first.
- Chooses that branch which allows the branching variable to take the value it has in the **current incumbent**.
- If this branch does not yield a solution then go to the other branch.



## EXAMPLE OF GUIDED DIVE

## Implementation

- RINS Requires a **call** to the MIP solver to solve a smaller sized MIP.
- In MINTO, this can be achieved through **recursion**.
- A new problem will be formulated using the APPL functions.
- The original formulation will be kept. The **branching constraints** will be dropped.
- Solve this simpler MIP.
- Return the result back to the parent.
- Based on the results obtained from the child, update the parent.
- In case of SYMPHONY, recursion is now possible.

## Implementation-II

- Local Branching:
  - No new MIP instance required.
  - Add soft bound constraints at a new incumbent and resolve.
  - May need to backtrack by removing the bound constraints.
  - Other issues:
    - \* Finding suitable parameters like maximum sub-problem size, branching parameters, time-limits etc.